## ICS 321 Fall 2011 Normal Forms 1

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## The Problem with Redundancy

Hourly_Emps

| SSN | Name | Lot | Rating | Hourly_wages | Hours_worked |
| :--- | :--- | :--- | :--- | :--- | :--- |
| 123-22-2366 | Attishoo | 48 | 8 | 10 | 40 |
| 231-31-5368 | Smiley | 22 | 8 | 10 | 30 |
| $131-24-3650$ | Smethurst | 35 | 5 | 7 | 30 |
| $434-26-3751$ | Guldu | 35 | 5 | 7 | 32 |
| $612-67-4134$ | Madayan | 35 | 8 | 10 | 40 |

- Suppose hourly wages are determined by rating
- Redundant storage : $(8,10)$ stored multiple times
- Update anomaly : change hourly wages in row 1
- Insertion anomaly : requires knowing hourly wages for the rating
- Deletion anomaly : deleting all $(8,10)$ loses info


## Using Two Smaller Tables

Hourly_Emps

| SSN | Name | Lot | Rating | Hours <br> worked |
| :--- | :--- | :--- | :--- | :--- |
| 123-22-2366 | Attishoo | 48 | 8 | 40 |
| 231-31-5368 | Smiley | 22 | 8 | 30 |
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| $612-67-4134$ | Madayan | 35 | 8 | 40 |

- Notation: denote relation schema by listing the attributes SNLRWH
- Update anomaly : Can we change W for Attishoo?
- Insertion anomaly : What if we want to insert an employee and don't know the hourly wage for his rating?
- Deletion anomaly : If we delete all employees with rating 5 , do we lose the information about the wage for rating 5 ?


## Decomposition

| Hourly_Emps |  |  |  |  |
| :--- | :--- | :--- | :--- | :--- |
| SSN | Name | Lot | Rating | Hours_ <br> worked |
| 123-22-2366 | Attishoo | 48 | 8 | 40 |
| 231-31-5368 | Smiley | 22 | 8 | 30 |
| 131-24-3650 | Smethurst | 35 | 5 | 30 |
| $434-26-3751$ | Guldu | 35 | 5 | 32 |
| $612-67-4134$ | Madayan | 35 | 8 | 40 |

RatingWages

| Rating | Hourly <br> wages |
| :--- | :--- |
| 5 | 7 |
| 8 | 10 |

- Remove redundancy by decomposition
- Since hourly wage is completely determined by rating, factor out hourly wage.
- Pros: less redundancy less anomalies
- Cons: retrieving the hourly wage of an employee requires a join


## Normal Forms

- Helps with the question: do we need to refine the schema?
- If a relation is in a certain normal form (BCNF, 3NF etc.), it is known that certain kinds of problems are avoided/minimized. This can be used to help us decide whether decomposing the relation will help.
- Role of FDs in detecting redundancy:
- Consider a relation R with 3 attributes, ABC .
- No FDs hold: There is no redundancy here.
- Given A $\rightarrow$ B: Several tuples could have the same A value, and if so, they'll all have the same B value!


## Boyce-Codd Normal Form (BCNF)

- Let $R$ denote a relation, $X$ a set of attributes from $R, A$ an attribute from $R$, and $F$ the set of FDs that hold over R.
- $R$ is in BCNF if for all $X \rightarrow A$ in $F^{+}$,
$-A \in X$ (trivial FD) or
-X is a superkey
The only non-trivial FDs that hold are key constraints
- Negation: R is not in BCNF if there exists an X $\rightarrow A$ in $F^{+}$, such that $A \notin X$ (non-trivial FD) AND $X$ is not a key


## Examples: BCNF

- Are the following in BCNF ?

| Firstname | Lastname | DOB | Address | Telephone |
| :--- | :--- | :--- | :--- | :--- |
| John | Smith | Sep 91979 | Honolulu,HI | 808-343-0809 |
| F= \{ FLD $\rightarrow$ FLDAT $\}$ |  |  |  |  |


| Firstname | Lastname | DOB | Street | CityState | Zipcode | Telephone |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- |
| John | Smith | Sep 9 <br> 1979 | 1680 East West <br> Rd. | Honolulu,HI | 96822 | $808-343-$ <br> 0809 |

$\mathrm{F}=\{\mathrm{FLD} \rightarrow \mathrm{FLDSCZT}, \mathrm{C} \rightarrow \mathrm{Z}\}$

## Third Normal Form (3NF)

- Let $R$ denote a relation, $X$ a set of attributes from $R, A$ an attribute from $R$, and $F$ the set of FDs that hold over R.
- $R$ is in 3NF if for all $X \rightarrow A$ in $F^{+}$,
$-A \in X$ (trivial FD) or
$-X$ is a superkey or
- A is part of some key
- Negation: $R$ is not in 3NF if there exists an $X \rightarrow A$ in $F^{+}$, such that $A \notin X$ (non-trivial FD) AND $X$ is not a key AND $A$ is not part of some key
- If $R$ is in BCNF, obviously in 3NF.
- If $R$ is in $3 N F$, some redundancy is possible. It is a compromise, used when BCNF not achievable (e.g., no "good" decomp, or performance considerations).


## Example: 3NF

- Which of the following is in 3NF and which in BCNF ?

| Firstname | Lastname | DOB |  | Address | Telephone |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| John | Smith | Sep 91979 |  | Honolulu, HI | 808-343-0809 |  |  |
| $\mathrm{F}=\{\mathrm{FLD} \rightarrow \mathrm{FLDAT}\}$ |  |  |  |  |  |  |  |
| Firstname | Lastname | DOB | Stre |  | CityState | Zipcode | Telephone |
| John | Smith | $\begin{aligned} & \text { Sep } 9 \\ & 1979 \end{aligned}$ |  | East West | Honolulu, HI | 96822 | $\begin{aligned} & \text { 808-343- } \\ & 0809 \end{aligned}$ |

$F=\{F L D \rightarrow F L D S C Z T, C \rightarrow Z\}$

| Student | Course | Instructor |
| :--- | :--- | :--- |
| Smith | OS | Mark |

$F=\{S C \rightarrow I, I \rightarrow C\}$

